

Connecting declarative software tools

Declarative tools [for] connecting software

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Summary

- Connecting declarative software tools:
 - ◆ The verifying compiler project
 - ◆ Concrete problems
 - ◆ Interoperability for declarative tools and languages
- Declarative tools for connecting software:
 - ◆ Models and logics for Web analysis and development
 - ◆ Declarative models for security protocols
- Conclusions and future work



Connecting declarative software tools



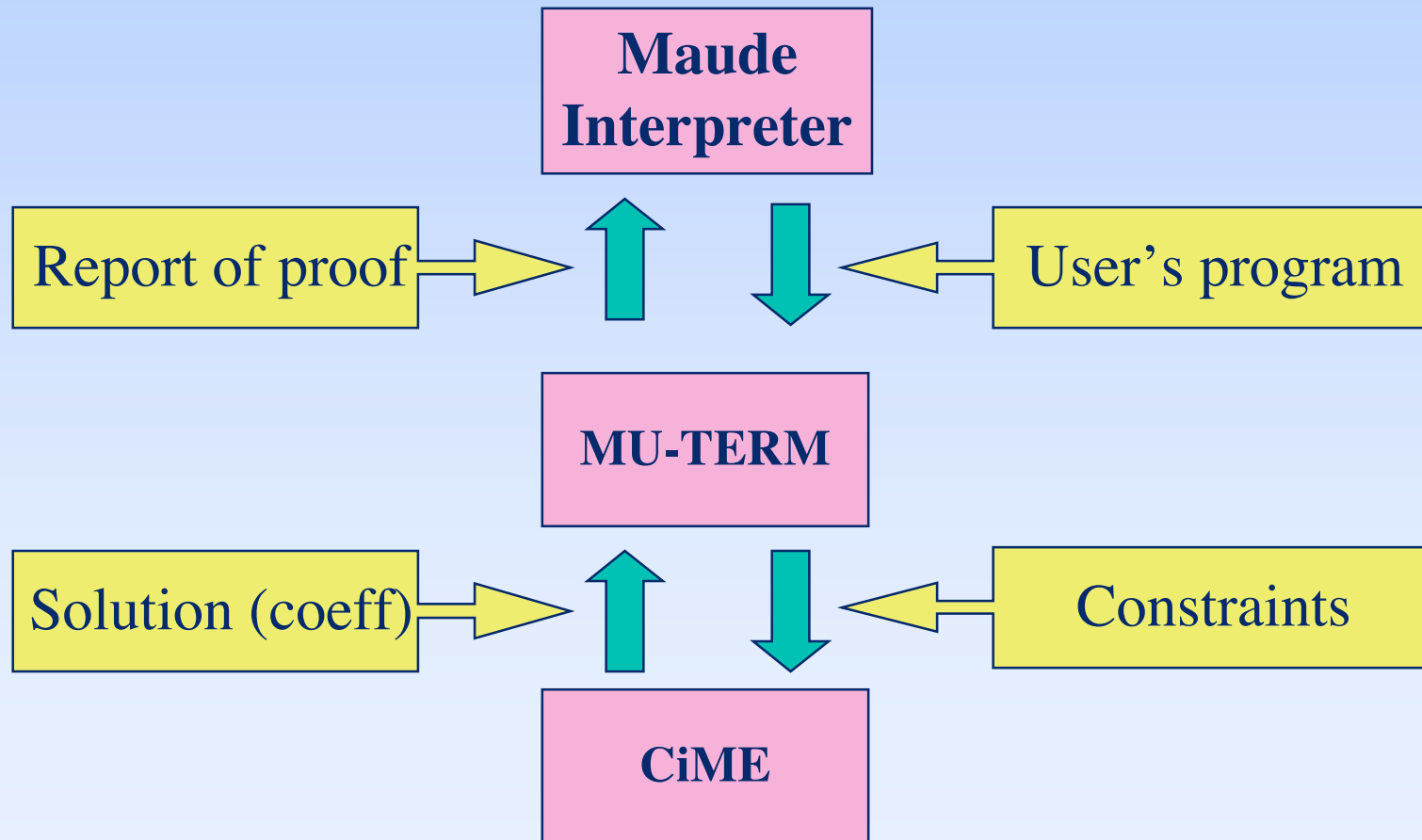
Connecting declarative tools

- As part of the 50th anniversary of the Journal of the ACM, an special issue of the journal by highly renowned researchers was published (*Journal of the ACM vol 50, issue 1, January 2003*)
- The aim was to establish the **most important challenges in Informatics and Computer Science for the XXI century**

Connecting declarative tools

- The verifying compiler: a grand (although classic!) challenge revisited by T. Hoare
- Program verification, program debugging, and program analysis will be essential components of such a tool
- Its effective development will require an **incremental** and **cooperative** effort from different work teams **all around the world**

Motivation: declarative languages



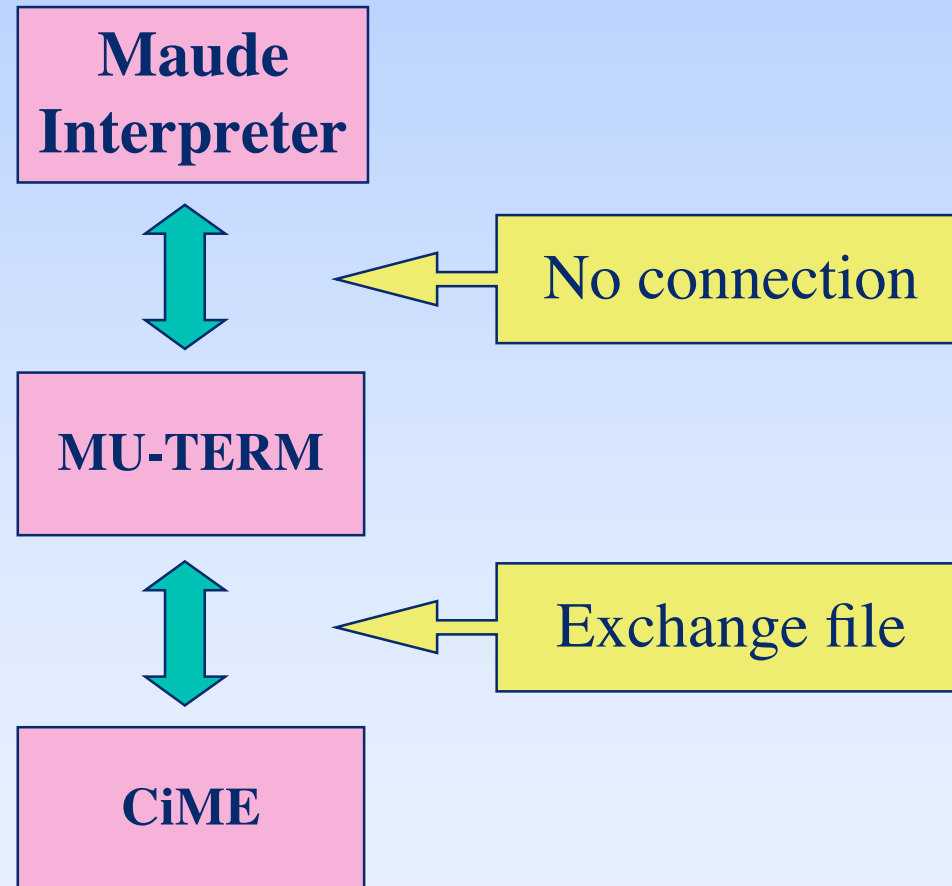
Motivation: declarative languages

*How to connect these tools for
automatically
proving termination of such programs?*



Connecting software tools: concrete problems

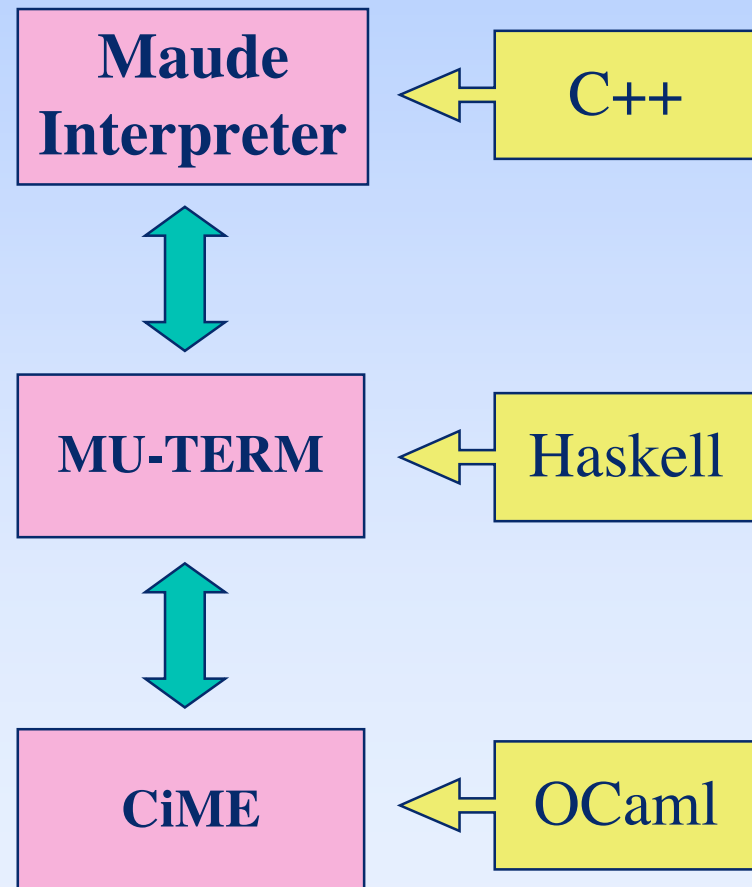
Connecting tools: concrete problems



Connecting tools: concrete problems

Data structures:

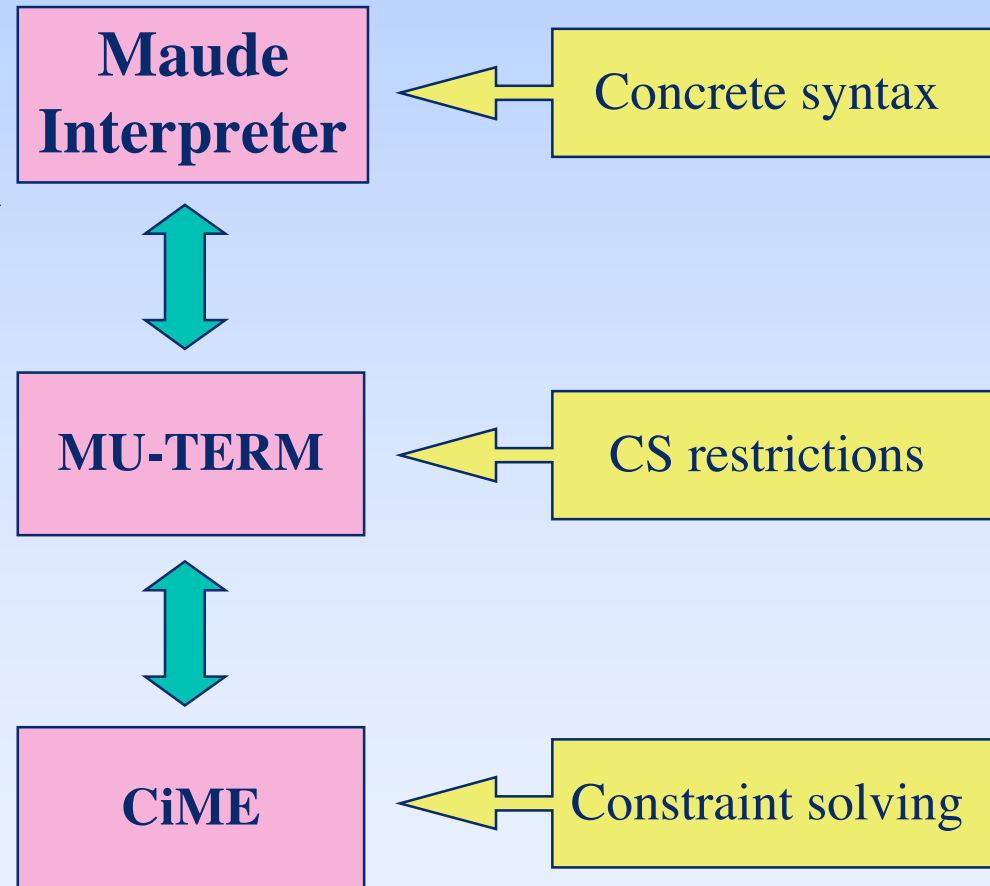
Although they could be linked as object modules, the data representations should be (made) compatible for exchanging data through primary memory



Connecting tools: concrete problems

Distributed:

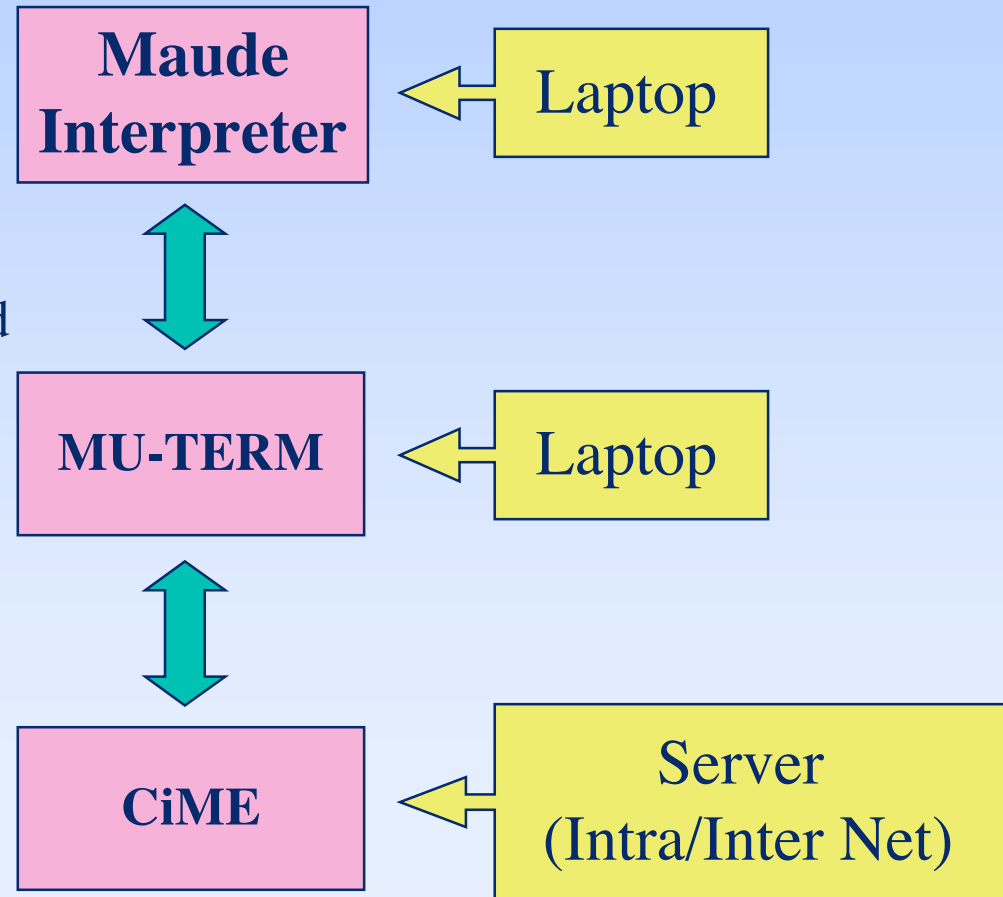
Proofs of termination of Programs involve different kinds of knowledge and expertise. Combining different tools to prove termination is often necessary



Connecting tools: concrete problems

Efficiency:

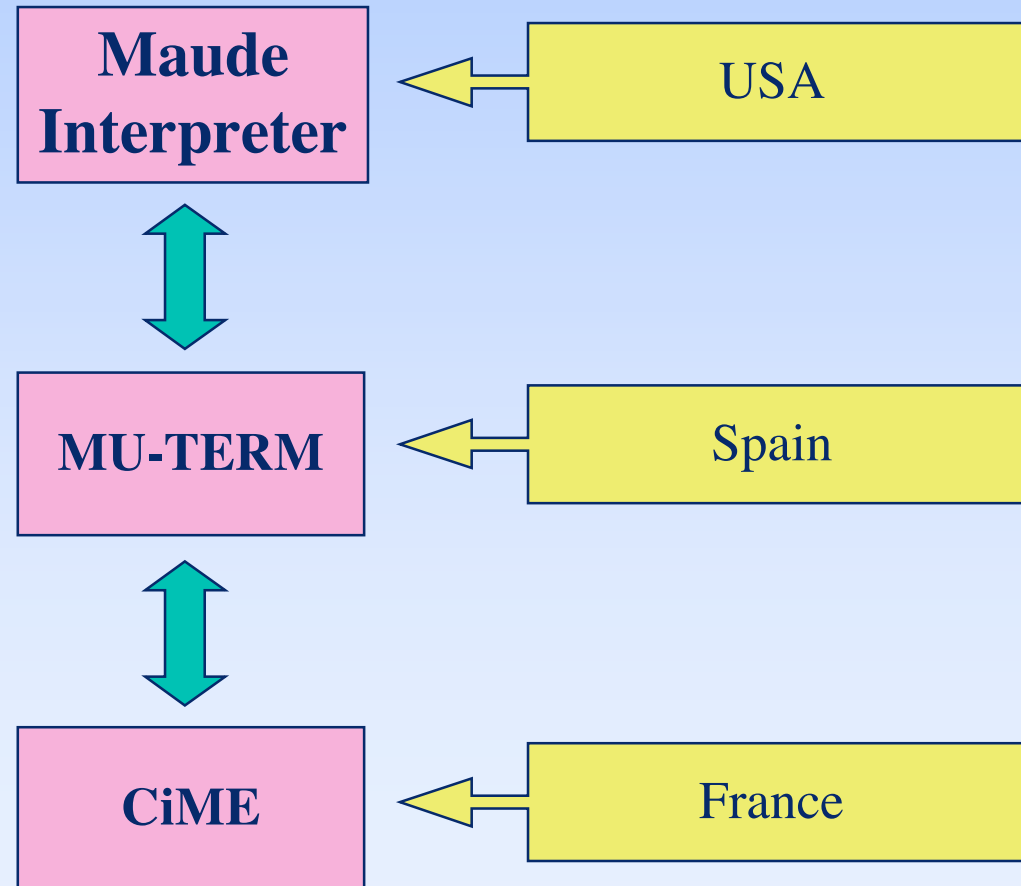
Proofs of termination involve search problems which are costly. Having specialized servers devoted to prove termination can be useful



Connecting tools: concrete problems

International:

Maude is developed and maintained (mainly) by the UIUC and SRI at USA;
MU-TERM has been made at the UPV (Spain)
CiME is being developed at the U. Paris VII (France)





Connecting applications: interoperability



Connecting applications

- Interoperability: making possible for a program on one system to get access to programs and data on another system
- Solutions: *Middleware systems*, e.g.,
 - ◆ COM
 - ◆ .NET
 - ◆ XML WWW Services

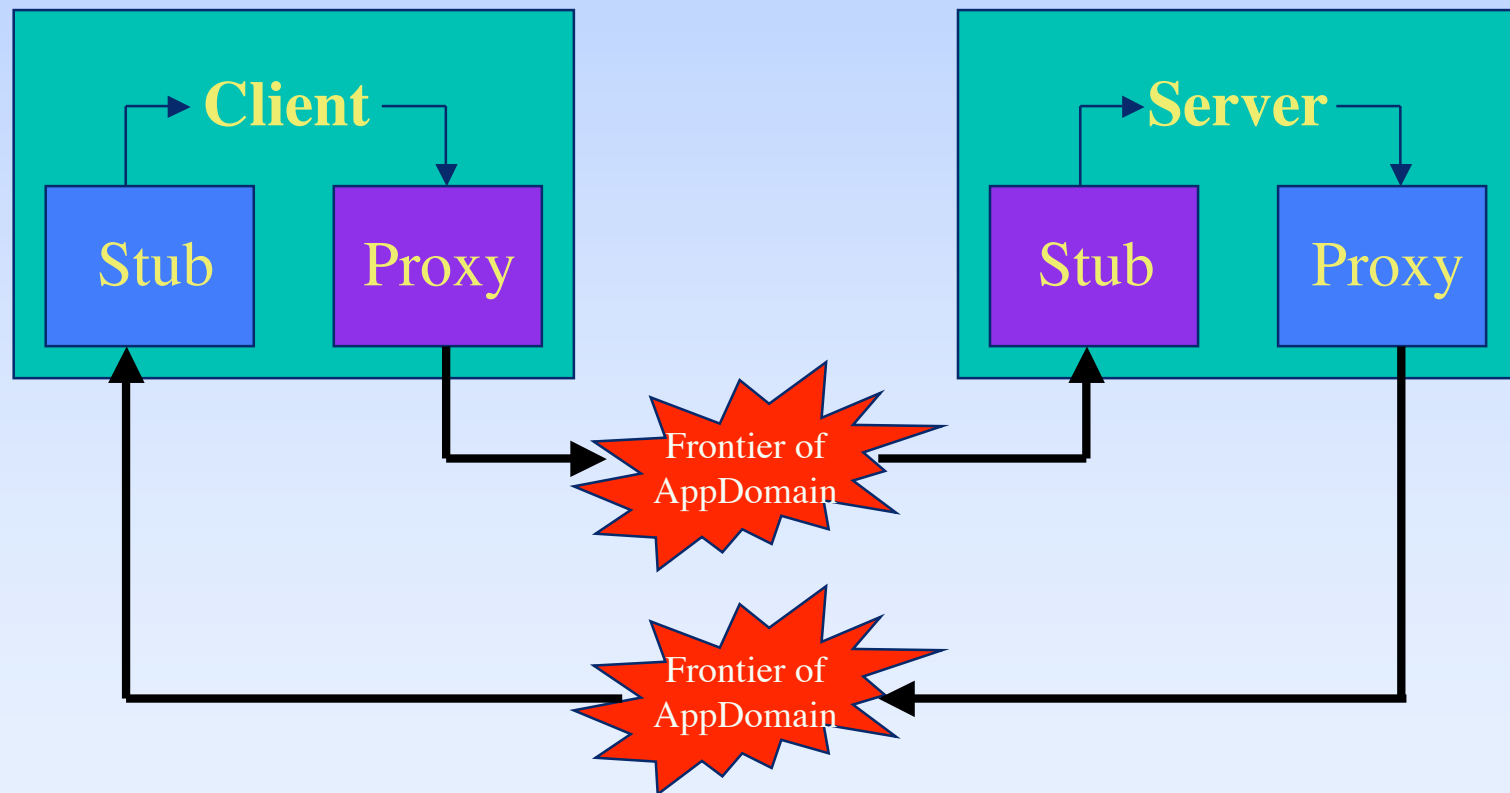
Connecting applications

■ Example: .NET:

- ◆ A core language (CLR) provides an abstract machine to implement more sophisticated languages:
 - ◆ C++ (or C#),
 - ◆ Java (or Java#)
 - ◆ ML,
 - ◆ Haskell (Mondrian), etc.
- ◆ The implementations can use a number of libraries (for GUIs, **remote access**,...)

Connecting applications

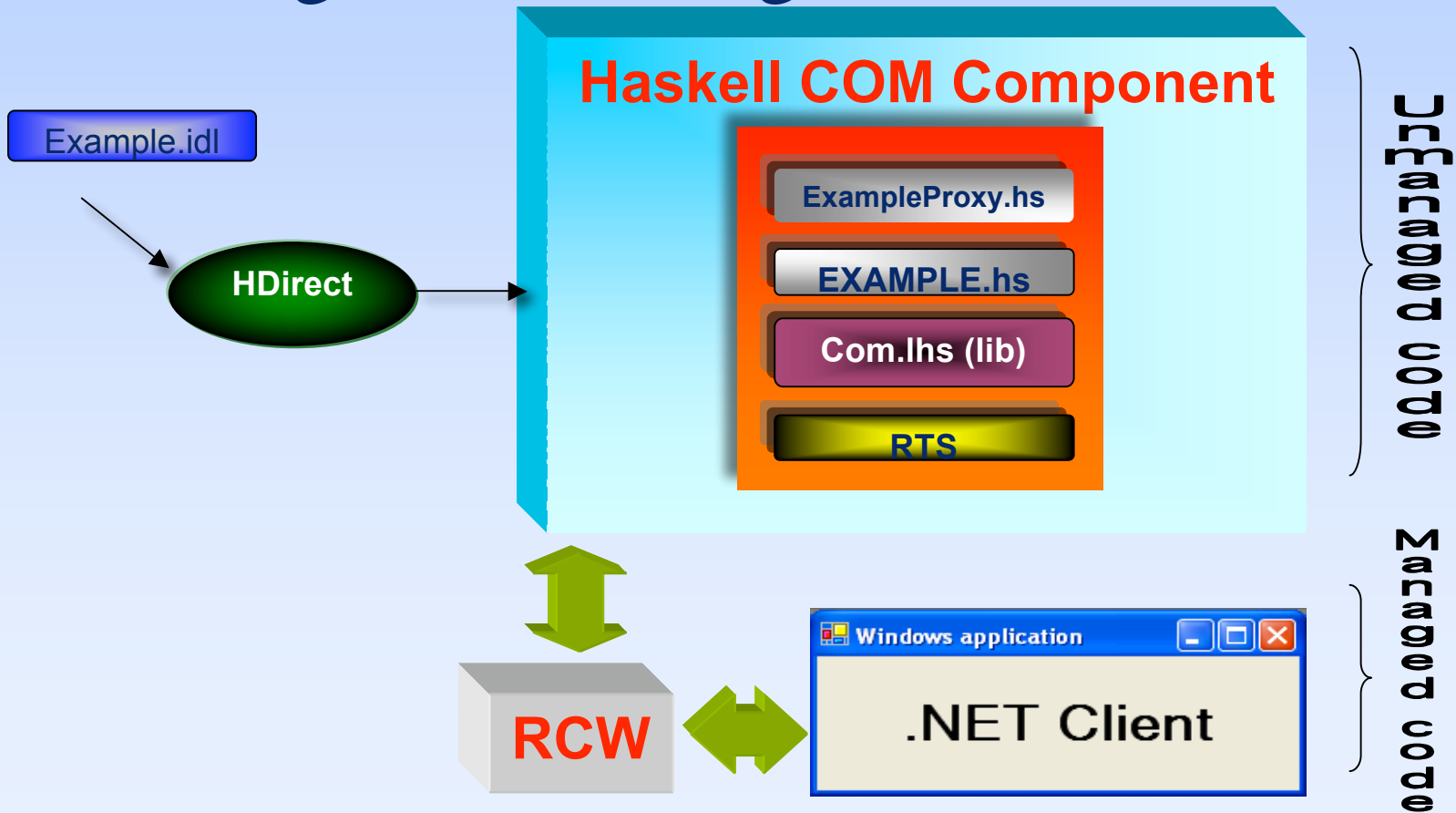
■ .NET Remoting:



AppDomains represent **local** or **remote** applications

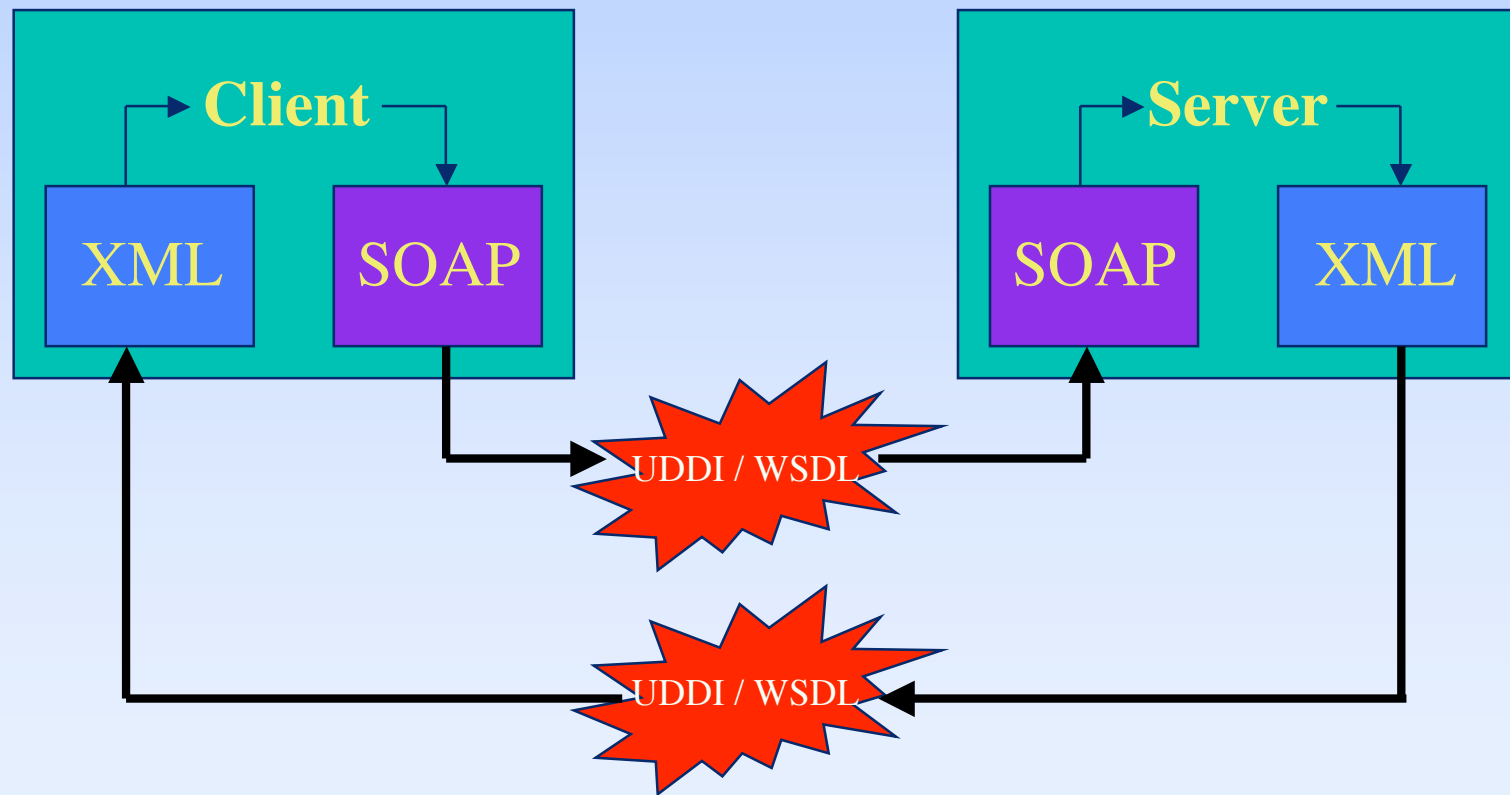
Connecting applications

■ Joining .NET through COM:



Connecting applications

■ WWW services:





Connecting applications

■ Common problems

- ◆ Exchanging data
- ◆ Defining remote services
- ◆ Finding external applications / servers
- ◆ Implementing remote calls
- ◆ Receiving results of remote calls



Connecting software tools: concrete actions

Connecting applications: actions

■ TPDB

- ◆ Recent common format for TRSs and termination problems:
 - ◆ Conditional equations / rules
 - ◆ Strategies
 - ◆ Type of problem (TRS, SRS, LP, ...)

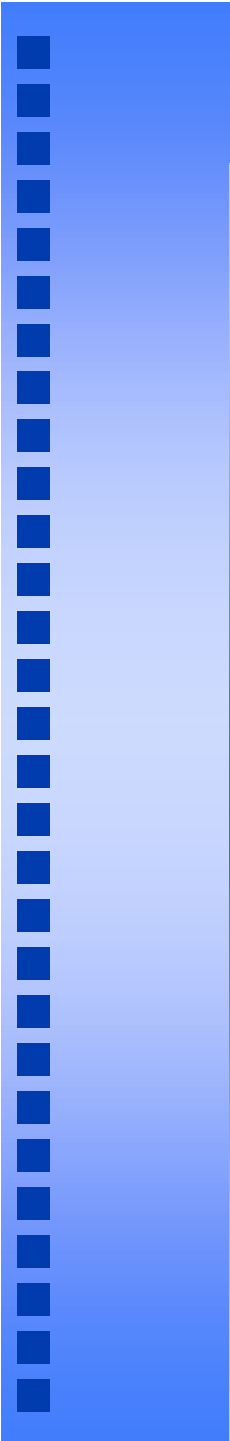
Connecting applications: actions

- Add information for specifying proofs
 - ◆ Simple / C_ε / DP-Simple termination
 - ◆ Constraint solving
 - ◆ Modular structure
 - ◆ Heuristics (and its combinations)
 - ◆ Ad-hoc partial / external proofs
- Use of XML for producing input / output information on proofs (e.g., for certification purposes)

Connecting applications: actions

This is an ambitious project which should eventually be agreed / addressed by the interested community.

Coordination with some technical groups (e.g., IFIP WG 1.6 or 1.3,...) would be interesting / desirable



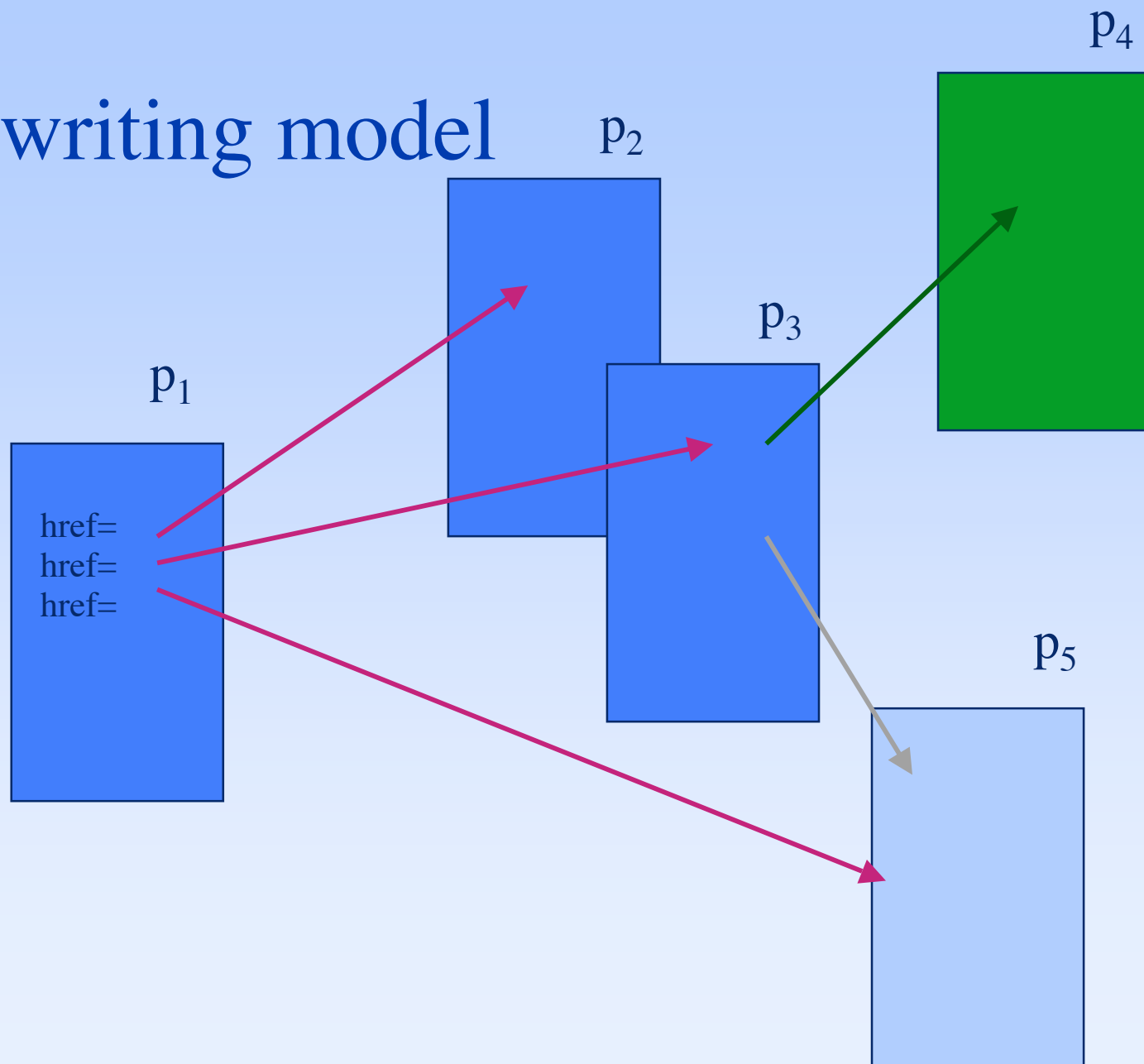
Declarative tools for connecting software



Declarative tools for connectivity

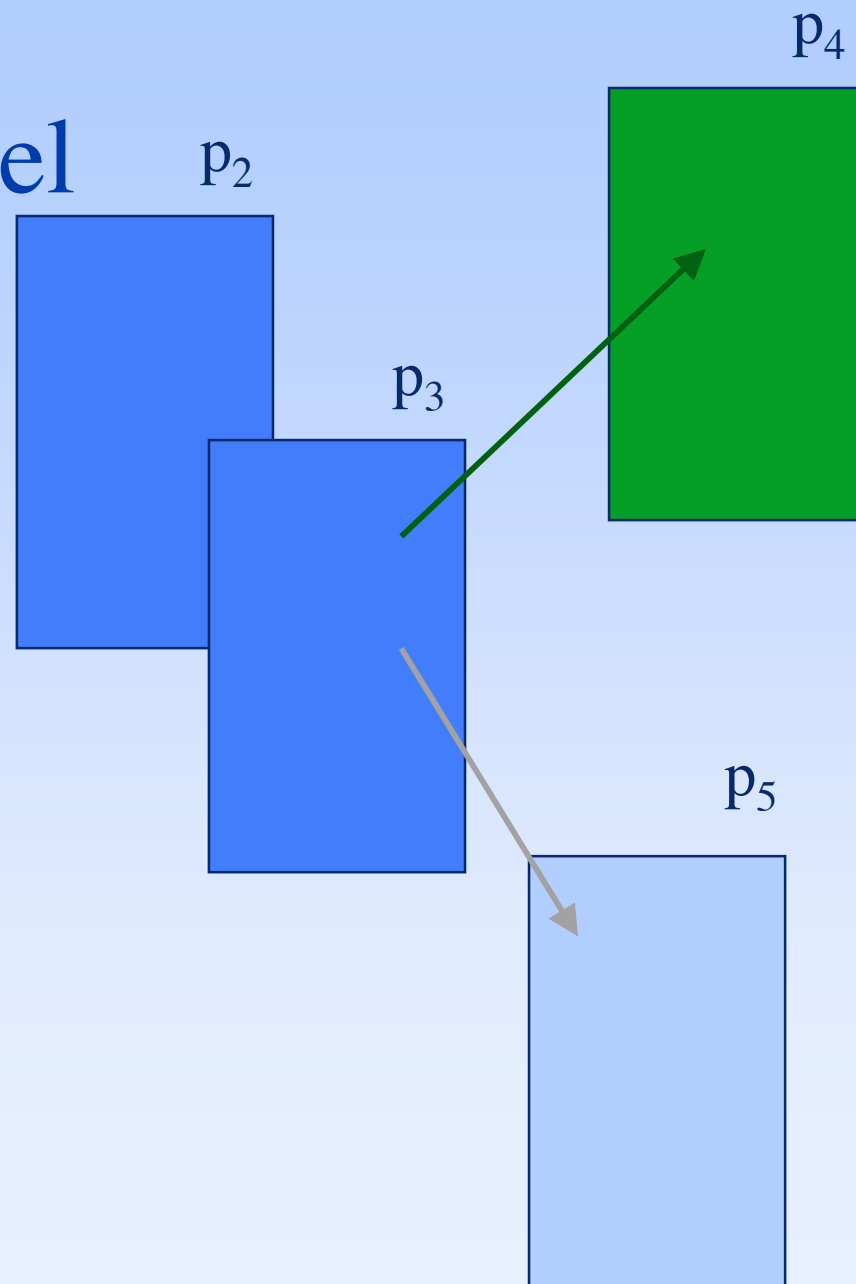
- Web site: a collection of connected Web pages
- Dynamic modeling: focus on the transitions between Web pages

Rewriting model



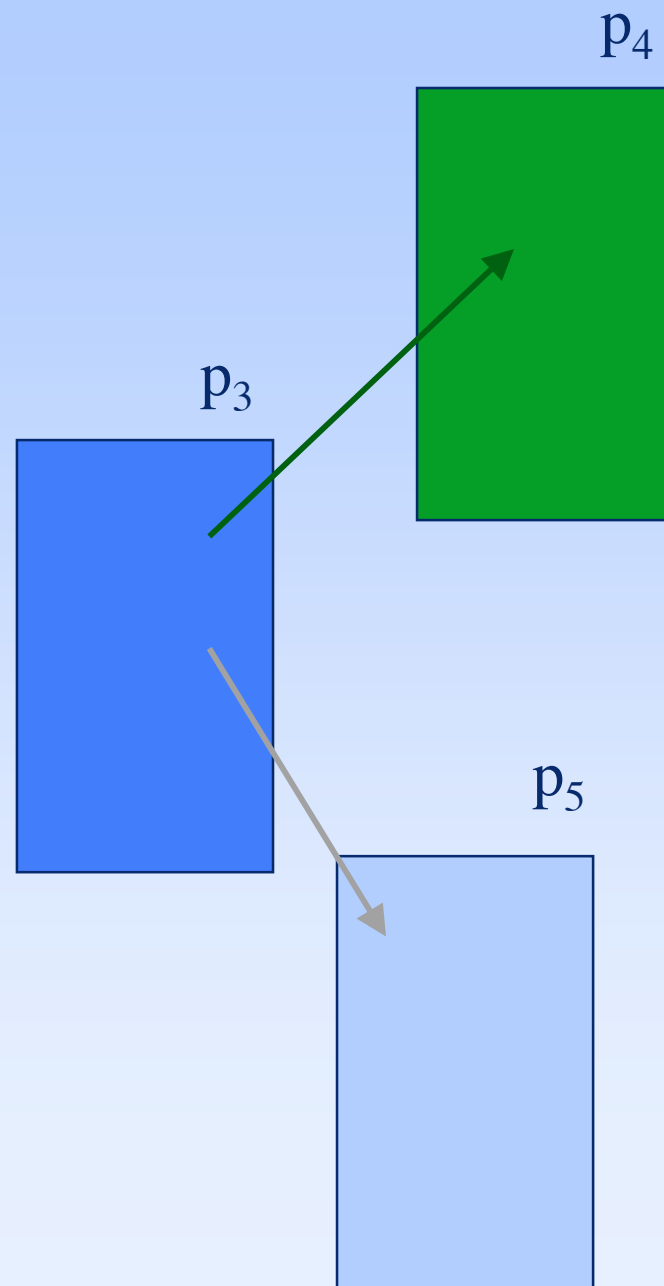
Rewriting model

$p_1(U) \rightarrow p_2(U)$
 $p_1(U) \rightarrow p_3(U)$
 $p_1(U) \rightarrow p_5(U)$



Rewriting model

$p_1(U) \rightarrow p_2(U)$
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Rewriting model

$p_1(U) \rightarrow p_2(U)$
 $p_1(U) \rightarrow p_3(U)$
 $p_1(U) \rightarrow p_5(U)$

$p_3(\mathbf{u}) \rightarrow p_4(\mathbf{u})$
 $p_3(\mathbf{u}') \rightarrow p_5(\mathbf{u}')$

p_4



p_5



Rewriting model

■ Term Rewriting System (TRS):

| | |
|-----------------------------|---|
| $p_1(U) \rightarrow p_2(U)$ | $p_3(\mathbf{u}) \rightarrow p_4(\mathbf{u})$ |
| $p_1(U) \rightarrow p_3(U)$ | $p_3(\mathbf{u}') \rightarrow p_5(\mathbf{u}')$ |
| $p_1(U) \rightarrow p_5(U)$ | |

■ Rewriting theories: first order logic (with variables ranging on terms) together with a binary predicate $R(x,y)$ associated to a TRS R :

- ◆ $R(x,y) = x \rightarrow y$: one-step rewriting theory
- ◆ $R(x,y) = x \rightarrow^* y$: rewriting theory

Rewriting model and logics

- **Example:** there is no ‘disconnected’ page:

$$\text{TM } y \exists x ((x \neq y) \wedge ((x \rightarrow y) \vee (y \rightarrow x)))$$

where ‘=’ is the predicate $R(x,y)$ associated to the empty TRS

- **Example:** there is no unreachable page (from the ‘main’ page):

$$\begin{aligned} &\text{TM } x (main \rightarrow^* x) \\ &\text{TM } x \exists u (main(u) \rightarrow^* x) \end{aligned}$$

Rewriting model and logics

- **Example:** “*there is no ‘disconnected’ page*”:

$${}^{\text{TM}}y \exists x ((x \neq y) \wedge ((x \rightarrow y) \vee (y \rightarrow x)))$$

where ‘=’ is the predicate $R(x,y)$ associated to the empty TRS

- **Example:** “*there is no unreachable page (from the ‘main’ page)*”:

$${}^{\text{TM}}x (main \rightarrow^* x)$$
$${}^{\text{TM}}x \exists u (main(u) \rightarrow^* x)$$

$${}^{\text{TM}}x (main(u_1) \rightarrow^* x) \vee \dots \vee (main(u_n) \rightarrow^* x)$$

Rewriting model: improvements

- **Example:** “no ‘unsafe’ access is possible”:

$${}^{\text{TM}}p \ {}^{\text{TM}}q \ {}^{\text{TM}}u \ {}^{\text{TM}}v \ ((p(u) \rightarrow^* q(v)) \Rightarrow (u=v))$$

- This is a **higher-order** sentence which does not belong to any rewriting theory!

Rewriting model: improvements

- This can be solved by introducing a new binary symbol to put together web pages and users as constant symbols: e.g., $\text{browse}(p,u)$

$$\text{TM } p \text{ TM } q \text{ TM } u \text{ TM } v ((\text{browse}(p,u) \rightarrow * \text{browse}(q,v)) \Rightarrow (u=v))$$

- Problem: no decidability results are available!!



Rewriting model: in practice

- Rewriting-based specification languages like **Maude** are well-suited to express dynamic models of Web sites
- In **Maude** a small query language is available (see the proceedings for some examples)
- Some existential queries are even possible on the basis of traversing the (finite) state space by using a breadth-first search strategy



Rewriting model: network protocols

- The **NRL Protocol Analyzer** (NPA) is a well-known tool for the formal specification and analysis of cryptographic protocols
- For the first time a precise formal specification of its *grammar-based techniques for invariant generation*, one of the main features of the NPA inference system, has been given
- This formal specification is given within the well-known framework of the **rewriting logic**



Conclusions / future work



Conclusions

- We are approaching the use of software tools with more complex systems (e.g., interpreters of programming languages)
- The combination of different tools with different expertise domain is required here



Conclusions

- Interoperability issues should be systematically considered when developing termination tools
- Rewriting-based logics are useful to model and analyze network systems and Web sites



Future work

- Which are the appropriate (**fragments of logics** which are useful to specify (and reason about) the dynamic behavior of Web sites?
- How **types, strategies, conditions**, etc. can help to get a more expressive model or to improve its power from a logic point of view (e.g., recovering decidability of the theories)



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